

MODEL ANSWER

SUMMER- 19 EXAMINATION

Subject Title: Digital Communication

Subject Code: 17535

Important Instructions to examiners:

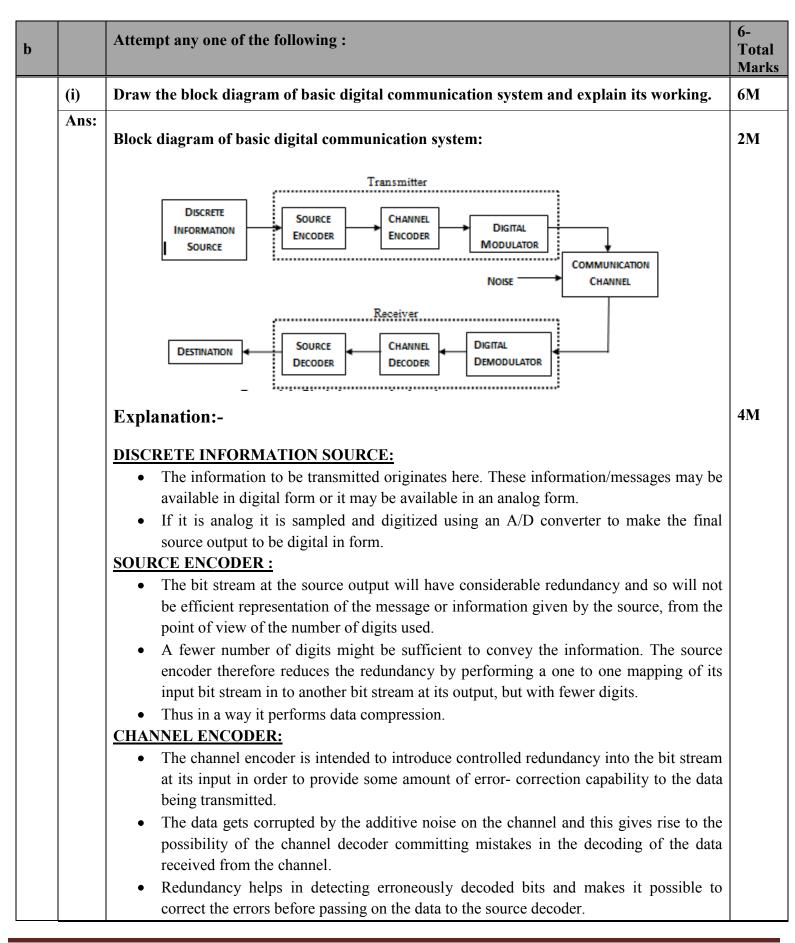
- 1) The answers should be examined by key words and not as word-to-word as given in the model answer scheme.
- 2) The model answer and the answer written by candidate may vary but the examiner may try to assess the understanding level of the candidate.
- 3) The language errors such as grammatical, spelling errors should not be given more Importance (Not applicable for subject English and Communication Skills.
- 4) While assessing figures, examiner may give credit for principal components indicated in the figure. The figures drawn by candidate and model answer may vary. The examiner may give credit for anyequivalent figure drawn.
- 5) Credits may be given step wise for numerical problems. In some cases, the assumed constant values may vary and there may be some difference in the candidate's answers and model answer.
- 6) In case of some questions credit may be given by judgement on part of examiner of relevant answer based on candidate's understanding.
- 7) For programming language papers, credit may be given to any other program based on equivalent concept.

	ii)	State sampling theorem and list its types.	4M
		Within the frequency limits of this channel i.e. over the bandwidth.	
		S/N – ratio of total signal power to total random noise power at the input of the receiver	
		C - Channel capacity bits/sec	
		S – Average signal power N – Average noise power within the channel bandwidth	
		where B – channel bandwidth (Hz)	
		$C = B \log 2 [1 + S/N] \text{ bits/sec}$	
		noise is given by:-	
		Channel capacity of a channel with bandwidth B and band limited additive Gaussian white	
		Mathematical expression:-	
		arbitrary small probability of error. The unit of channel capacity is <i>bits/sec</i> .	
		It is defined as the maximum rate at which data can be transferred over the channel with an	
		Channel capacity:-	
	Ans:	(Definition - 2M, Mathematical expression -2M)	
	i)	Define channel capacity with mathematical expression.	4M
		Define shownel consists with mothematical symposium	Marks
Q.1		Attempt any THREE of the following :	Total
_			e 12-
No.	Q.N.		Schem
Q.	Sub	Answer	ng
			Marki



 SAMPLING THEOREM: Sampling theorem states that a band-limited signal of finite frequency component <i>fm</i> Hz can be represented and recove Samples taken at a rate of <i>fs</i> samples per second provi sampling Frequency. Types of Sampling:- Natural sampling Flat top Sampling:- Natural sampling Flat top Sampling:-		
 Ans: Need of Multiplexing:- Multiplexing is the set of techniques that allows the sir signals across a signals data link. When many signals or channels are to be transmitted sends simultaneously i.e. multiplexer converts many ir also all input we get simultaneously. Sending many signals separately is expensive and requa a need of multiplexing. Efficient utilization of channel capacity. For example in cable T.V distributor sends many channel capacity. Ket application of spread specture modulation. Ans: (Each application - 1M) Applications Of Spread Specture Modulation:- The Spread Specture Communications are widely use Avionics, Scientific, and Civil uses. The applications include Jam-resistant communication systems CDMA radios High Resolution Ranging: Spread Specture Conresolution ranging. It is possible to locate an of techniques. for example where it could be use (GPS). WLAN: Wireless LAN (Local Area Networ communications.	red completely from a set of	
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 ii. Direct Sequence Spread Spectrum Com iii. Frequency Hopping Spread Spectrum C 5. Cordless Phones 6. Long-range wireless phones for home and industr 7. Cellular base stations interconnection. 	he the following: nmunications is often used in high bject with good accuracy using SS sed is Global Positioning System ks) widely use spread spectrum	4M

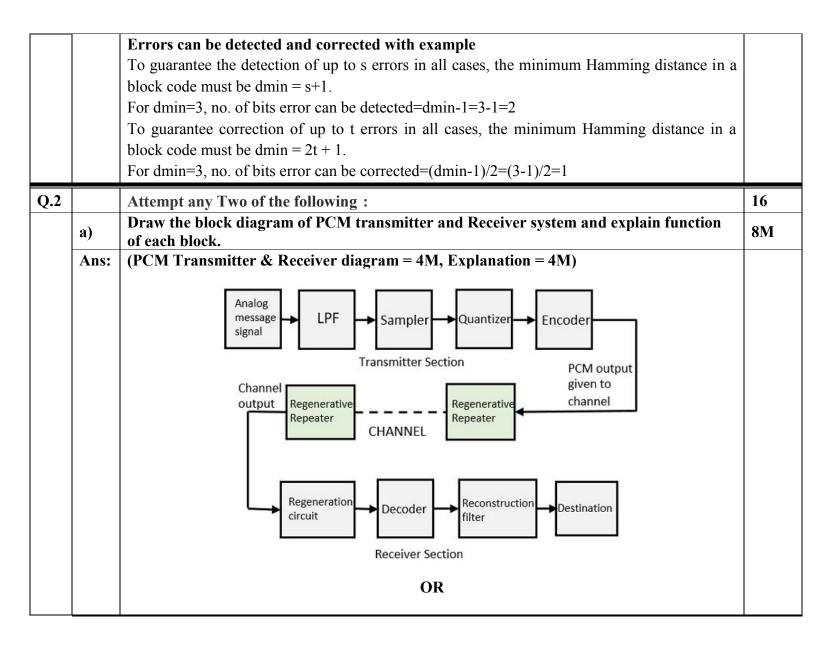




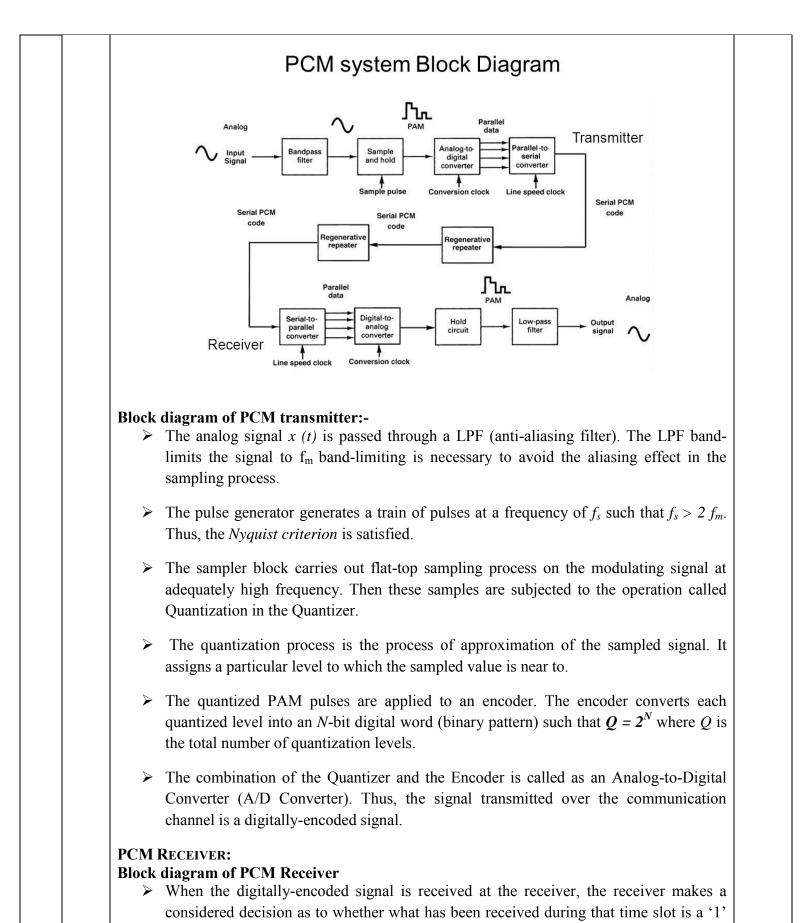


 DIGITAL MODULATOR: The physical channels are basically analog in nature; the digital modulator takes each digital binary digit at its input and maps it, in a one -to - one fashion, into a continuous waveform. PHYSICAL CHANNEL: The digitally modulated signal is passed on to the physical channel, which is nothing 	
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continuous waveform. <u>PHYSICAL CHANNEL:</u>	
PHYSICAL CHANNEL:	
• The digitally modulated signal is passed on to the physical channel, which is nothing	
but the physical medium through which the signals are transmitted.	
• It may take a variety of forms- a pair of twisted wires, coaxial cable, a wave guide, a microwave radio, or an optical fiber.	
 During its passage through the channel, the signal gets corrupted by noise. This noise 	
may be thermal noise originating from electronic circuits or atmospheric noise, or	
manmade noise, or as is generally the case, a combination of most or all of them.	
THE DIGITAL DEMODULATOR:	
• The digital demodulator of the receiver receives the noise corrupted sequence of	
waveforms from the channel and by inverse mapping tries to give at its output, an	
estimate of the sequence of the binary digits that were available at the input of the	
digital modulator at the transmitting end.	
THE CHANNEL DECODER:	
• The output sequences of digits from the digital demodulator are fed to the channel	
decoder. Using its knowledge of the type of coding performed by the channel encoder	
at the transmitting end and using the redundancy introduced by the channel encoder, it	
produces as its output, the output of the source coder of the transmitter with as few	
errors as possible.	
errors as possible. THE SOURCE DECODER:	
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	or a '0' i.e. a 'pulse' or 'no pulse'.	
	 Decision is made during each time slot, preferably at the center of the time slot, compare the sample amplitude with a fixed pre-set threshold and declare it as a '1' if it exceeds the threshold and declare it as a '0' during that time slot, if it is less than the threshold. 	
	Once it is decided that it is a 1 in a time slot, local pulse generator in the receiver is triggered and it gives a clean, noise free rectangular pulse. Although the received pulses were distorted and corrupted by additive noise a clean pulse is locally generated this process is known as 'regeneration'.	
	This sequence of clean 'pulse' and 'no-pulse', obtained through the process of regeneration are then fed to a serial to parallel converter then to a decoder which converts each code word into the corresponding quantized sample value.	
	These samples are then passed through a low pass reconstruction filter, which reconstructs an analog signal from these samples. This analog signal will be an approximation to the original analog message signal that was transmitted.	
b)	Illustrate BFSK signal generation with block diagram and waveforms. State bandwidth requirement and draw its frequency spectrum.	8M
	Spectrum = 2M)	
	Block diagram of BFSK signal Generation:-	
	Block diagram of BFSK signal Generation:- INPUT DATA (UNIPOLAR NRZ) b(t) $\sqrt{2P_s}\cos\omega_1 t$ $\sqrt{2P_s}\cos\omega_1 t$ $\sqrt{2P_s}\cos\omega_1 t$ INVERTER $\sqrt{2P_s}\cos\omega_1 t$ $\sqrt{2P_s}\cos\omega_1 t$	
	INPUT DATA (UNIPOLAR NRZ) b(t) $\sqrt{2P_s \cos \omega_1 t}$ $\sqrt{2P_s \cos \omega_1 t}$	
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FSK TRANSMITTER:

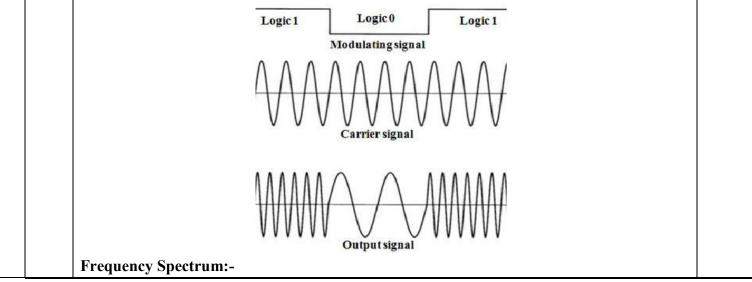
• In general, for binary FSK, the carriers can be represented by:

For binary 0, $V_0(t) = \sqrt{2P_s} \cos 2\pi f_0 t = \sqrt{2P_s} \cos \omega_0 t$

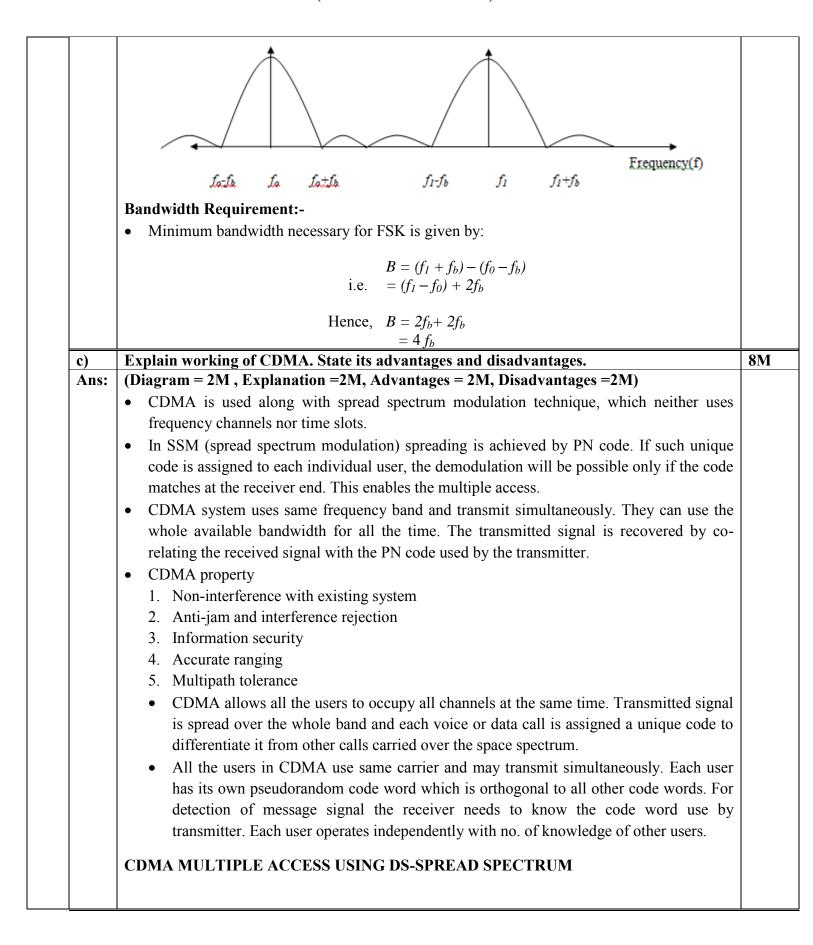
For binary 1, $V_l(t) = \sqrt{2P_s} \cos 2\pi f_l t = \sqrt{2P_s} \cos \omega_1 t$

- Assume that high frequency is transmitted for 1 and low frequency is transmitted for 0
- Low frequency = $f_0 = f_c \delta f$
- High frequency = $f_1 = f_c + \delta f$
- The two carriers may be generated from separate oscillators independent of one another as shown in the above figure.
- As shown in figure the input binary data is given directly to the multiplier and is inverted and given to second multiplier.
- Two different carriers have different frequency generated by the two oscillators and applied to the multipliers.
- The output of both the multipliers is an ASK signal which is added by the summer. Thus, the output of the adder is the FSK wave.
- Since there are two different oscillators used for the carrier signal the combined signal at the summer therefore has discontinuous amplitude and phase which is undesirable in FSK. Hence, we need a common carrier for FSK modulation to avoid these discontinuities in amplitude and phase.

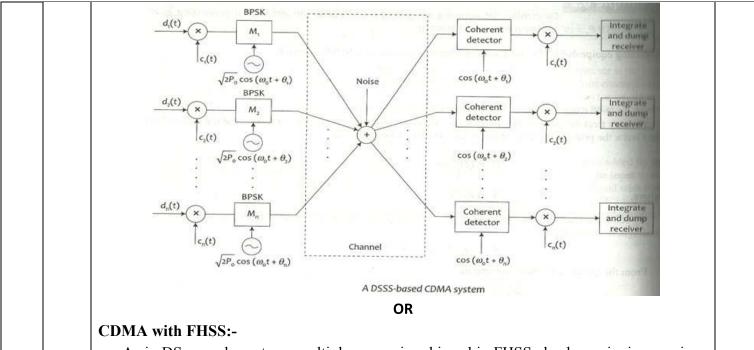
Waveform:-



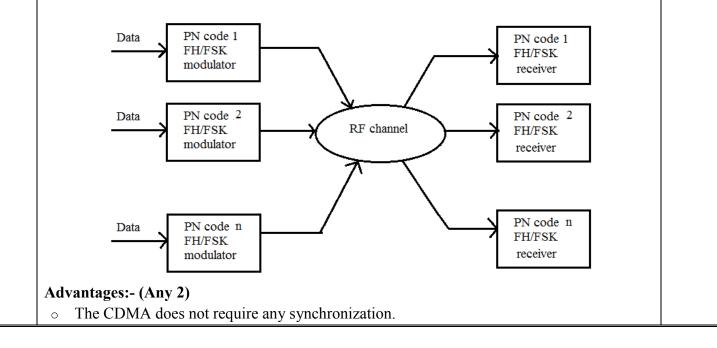




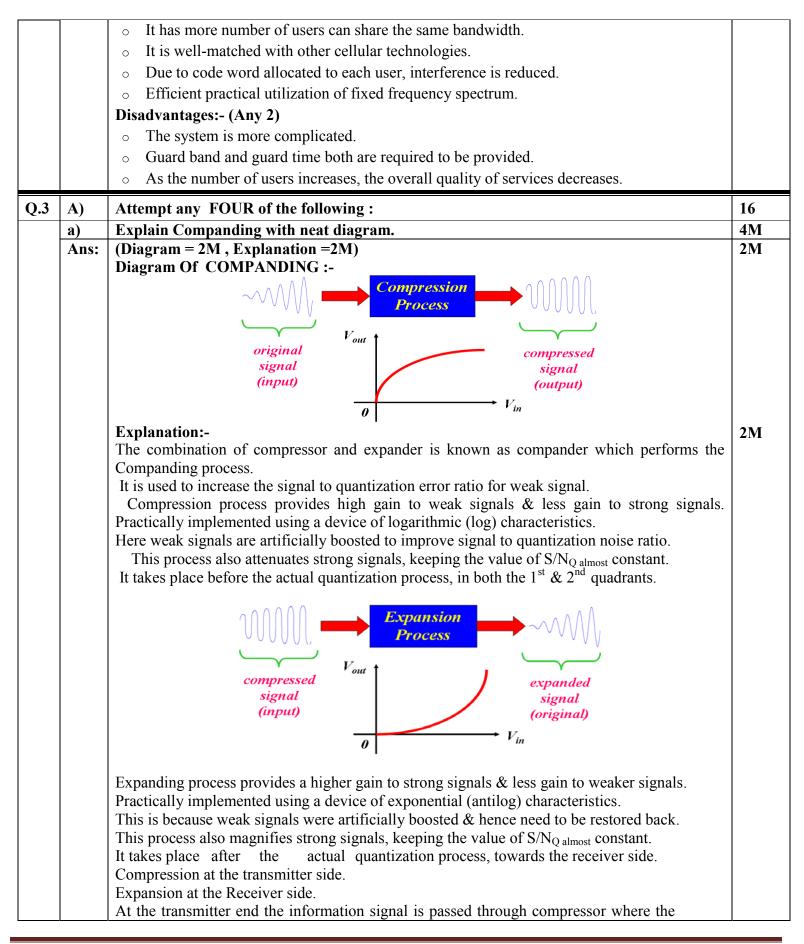




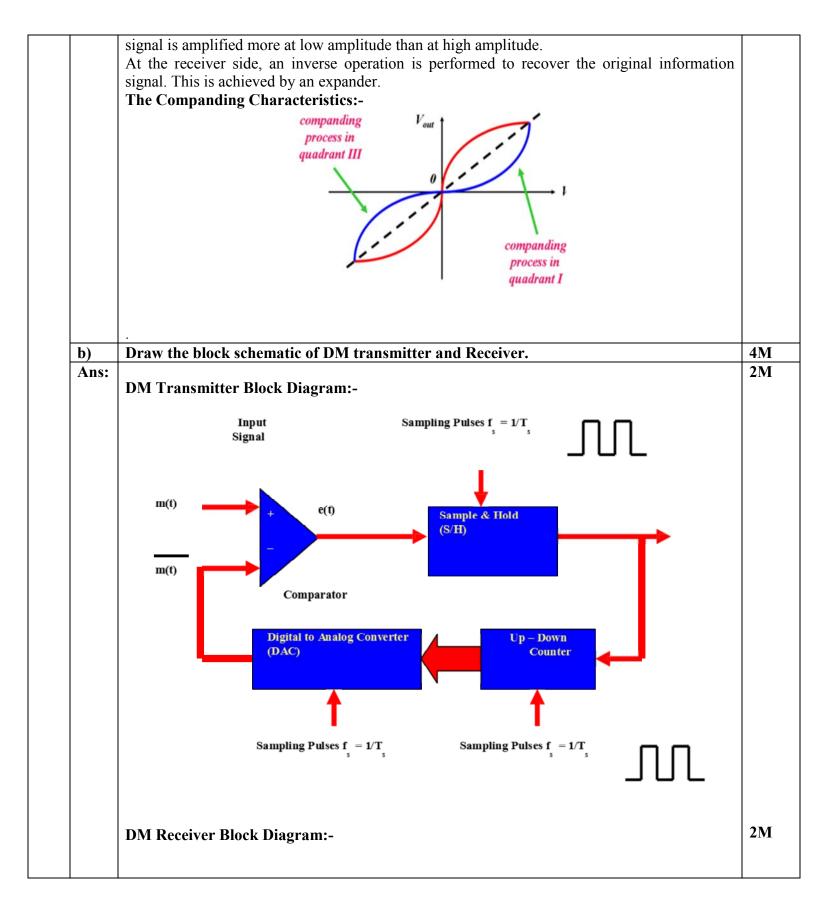
- As in DS spread spectrum multiple access is achieved in FHSS also by assigning a unique PN code to each user, which in this case controls the frequency-hopping pattern.
- These codes that are assigned must be so chosen that collision do not occur.
- The frequency produced by the frequency synthesizer during a chip period depends on the PN sequence value during that chip period.
- So sometimes it may so happen that two or more users have at a given time, the same PN sequence values produced by their respective PN sequence generators. In that case a collision is said to have occurred in the spectrum.
- Whether it is slow hopping or fast hopping FHSS, when a collision occurs it results in considerable increase in detection errors.



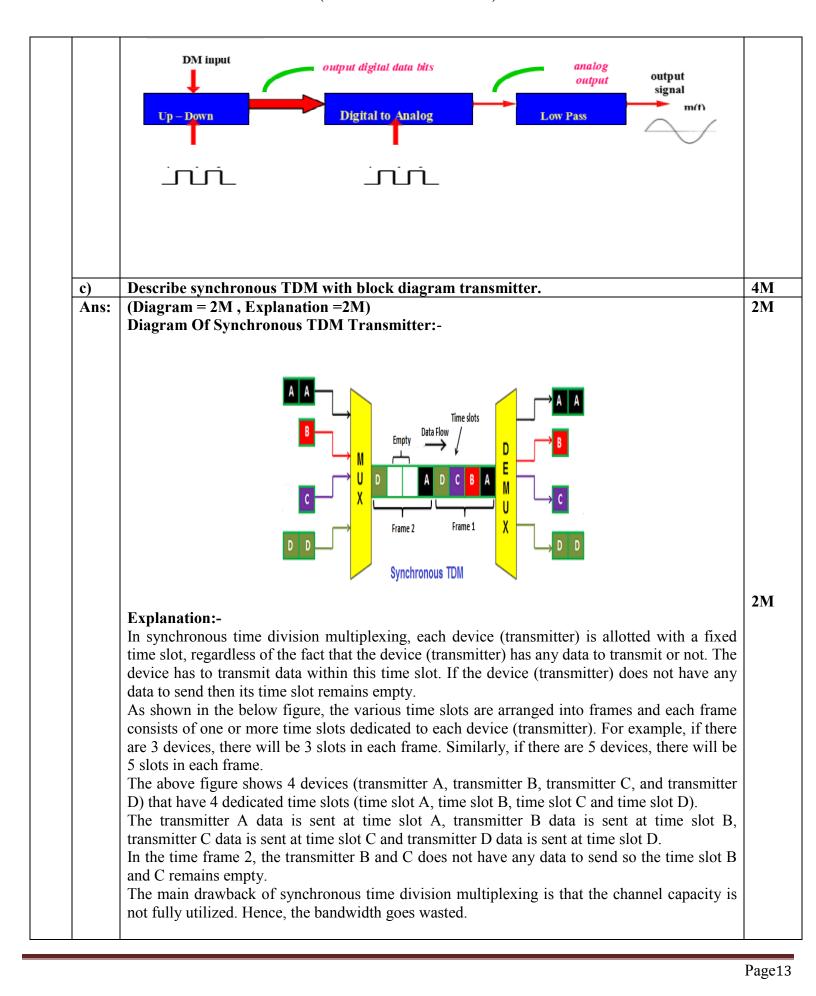














,	Compare	QPSK QASK (any four 			4]
Ans:		Parameter	QPSK	QASK	E
		1.Information is	phase	Amplitude and phase	P
		transmitted by change			1
		in			
		2. Number of	N=2	N=3 Or 4 Or 5 & so on	
		bits/symbol			
		3.Number of possible	Four	M=2 ^N	
		symbols			
		4.Detection method	coherent	coherent	
		5.minimum bandwidth	fb	2fb/N	
		6.symbol duration	2Tb	NTb	
		7.Noise immunity	Comparatively	Better than QPSK	
			less than QAM	Better than QI SK	
		8.System complexity		More complex then	
		8.System complexity	Less complex than	More complex than	
		0.0.1.1.1.	QAM	QPSK	
		9.Probability of error	More than QAM	Less than QPSK	
		10.Performance of	Less than QAM	Better than QPSK	
		system			
e)		QAM transmitter with bl			4]
Ans:	QAM tran	smitter block diagram:-			2
		Binary information		$ADDER \longrightarrow V_{QAM}(t)$ $A_{o}(t)\sqrt{2P_{s}} \sin \omega_{s} t$	
				4	
		CIOCK	$\sqrt{2P_s} \sin \omega_c$	^c t	
	-	smitter Explanation:-			1
	The bi stre	smitter Explanation:- am b _{(t}) is applied to the s	erial to parallel conve	erter, operating on a clock	which has
	The bi stre period of 7	Explanation: am $b_{(t)}$ is applied to the so Γ_{s} , Which is the symbol due	erial to parallel conve		which has
	The bi stre period of 7 presented i	Explanation: cam $b_{(t)}$ is applied to the set Γ_{s} . Which is the symbol due n the parallel form.	erial to parallel conve uration. The bits $b(^{t})$	erter, operating on a clock	which has
	The bi stre period of 7 presented i The four bi	Explanation: am $b_{(t)}$ is applied to the set Γ_{s} . Which is the symbol due n the parallel form. it symbols are $b_k + 3$, $bk + 3$.	erial to parallel conve uration. The bits $b(^{t})$ 1, and bk.	erter, operating on a clock are stored by the converte	which has er and then
	The bi stre period of T presented i The four bi Out of thes	EXAMPLE 1 Second Structure Explanation: $\Gamma_{s_{1}}$ which is applied to the second structure $\Gamma_{s_{1}}$ which is the symbol due of the parallel form. It symbols are b_{k} +3, bk +3 F_{se} four bits the first two bits	erial to parallel conve uration. The bits b(^t) a 1, and bk. ts are applied to a D/A	erter, operating on a clock	which has er and then
	The bi stre period of T presented i The four bi Out of thes applied to	Explanation: arm $b_{(t)}$ is applied to the second D/A converter.	erial to parallel conve uration. The bits b(^t) a 1, and bk. ts are applied to a D/A	erter, operating on a clock are stored by the converte A converter and the other t	which has er and then wo bits are
	The bi stre period of T presented i The four bi Out of thes applied to The Outpu	Assmitter Explanation: arm $b_{(t)}$ is applied to the set $\Gamma_{s_{t}}$ Which is the symbol due in the parallel form. it symbols are $b_{k} + 3$, $bk + 3$ set four bits the first two bits the second D/A converter. t of the first converter is A	erial to parallel conve uration. The bits $b(^{t})$ and bk. ts are applied to a D/A $A_{e}(t)$,which is modulat	erter, operating on a clock are stored by the converte A converter and the other t ed by the carrier $\sqrt{2}$ Ps cos	which has er and then wo bits are
	The bi stre period of T presented i The four bi Out of thes applied to The Outpu T whereas	Explanation: am $b_{(t)}$ is applied to the set Γ_{s} . Which is the symbol due in the parallel form. it symbols are $b_k + 3$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 4$, $bk + 3$, $bk + 4$, $bk $	erial to parallel conve uration. The bits $b(^{t})$ and bk. ts are applied to a D/A $A_{e}(t)$,which is modulat	erter, operating on a clock are stored by the converte A converter and the other t	which has er and then wo bits are
	The bi stre period of T presented i The four bi Out of thes applied to The Outpu T whereas ∞t in the b	Assmitter Explanation: arm $b_{(t)}$ is applied to the set Γ_{s} . Which is the symbol due in the parallel form. it symbols are $b_k + 3$, $bk + 3$, bk	erial to parallel conve uration. The bits $b(^{t})$ and bk. ts are applied to a D/A $A_{e}(t)$, which is modulat and D/A converter A_{o} is	erter, operating on a clock are stored by the converte A converter and the other t ed by the carrier $\sqrt{2}$ Ps cos is modulated by the carrie	which has er and then wo bits are
	The bi streperiod of T presented i The four bi Out of these applied to The Outpu T whereas ∞t in the bi $A_e(t) A_o(t)$	EXAMPLE 1 Substituting the second D/A converter is A is the second D/A converter is A is the output of the second and the s	erial to parallel conve uration. The bits $b(^{t})$ a 1, and bk. ts are applied to a D/A $A_{e}(t)$,which is modulat and D/A converter A_{o} is ted by the converter -2	erter, operating on a clock are stored by the converte A converter and the other t ed by the carrier $\sqrt{2}$ Ps cos is modulated by the carrie	which has er and then wo bits are 5∞ er $\sqrt{2}$ Ps sin
	The bi streperiod of T presented i The four bi Out of these applied to The Outpu T whereas ∞t in the bi $A_e(t) A_o(t)$	Assmitter Explanation: Fram $b_{(t)}$ is applied to the set Γ_{s} . Which is the symbol due in the parallel form. It symbols are $b_k + 3$, $bk + 4$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$,	erial to parallel conve uration. The bits $b(^{t})$ and bk. ts are applied to a D/A $A_{e}(t)$,which is modulat d D/A converter A_{o} is ted by the converter -2 e added together to ga	erter, operating on a clock are stored by the converte A converter and the other t ed by the carrier $\sqrt{2}$ Ps cos is modulated by the carrie 3 -1, +1,+3 volts et the QASK output signa	which has er and then wo bits are 5∞ er $\sqrt{2}$ Ps sin
	The bi streperiod of T presented i The four bi Out of these applied to T The Outpu T whereas ∞t in the bi $A_e(t) A_o(t)$ The balance	Assmitter Explanation: Fram $b_{(t)}$ is applied to the set Γ_{s} . Which is the symbol due in the parallel form. It symbols are $b_k + 3$, $bk + 4$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$, $bk + 3$, $bk + 3$, $bk + 3$, $bk + 4$, $bk + 3$,	erial to parallel conve uration. The bits $b(^{t})$ a 1, and bk. ts are applied to a D/A $A_{e}(t)$,which is modulat and D/A converter A_{o} is ted by the converter -2	erter, operating on a clock are stored by the converte A converter and the other t ed by the carrier $\sqrt{2}$ Ps cos is modulated by the carrie 3 -1, +1,+3 volts et the QASK output signa	which has er and then wo bits are 5∞ er $\sqrt{2}$ Ps sin



1 a	a)	Attempt any THREE of the following:	12
((i)	State advantages and disadvantages of digital communication system.	4M
A	Ans:	(Any 4 advantages = 2M , Any 2 disadvantages =2M)	2M
		Advantages of digital communication System:	
		1. Digital systems are simple and easy to build.	
		2. Insensitive to variations in atmospheric conditions like temperature, humidity etc. so	
		highly robust.	
		3. Storage and retrieval of voice, data or video is easy& inexpensive.	
		4. It offers considerable flexibility as voice, data, video can all be multiplexed using TDM, signal and image processing operations like compression of voice and image signal can be easily used etc	
		5. Cost of digital communication systems are coming down because of improvement of VLSI technology and available of IC's at ever decreasing price.	
		6. Error correction codes ensure fairly good protection against noise and interferences.	
		7. Powerful encryption and decryption algorithms are available for digital data so as to maintain high level of secrecy of communication.	
		8. In digital communication system, repeaters are used as regenerators so signal reaching at the destination can be almost error free. So this system can be used for long-haul communication.	
		9. Multiple data can be send simultaneously using multiplexing.	2M
		Disadvantages of Digital Communication System:	
		1. The transmission of digitally encoded analog signals requires significantly more bandwidth.	
		2. Digital transmission requires precise time synchronization between the clocks in the transmitter and receiver.	
		 Digital transmission systems are incompatible with older analog transmission systems. Digital communication system requires greater bandwidth 	
	(ii)		4M
-		Explain the need of using adaptive delta modulation.	4M
P	Ans:	Necessity of adaptive delta modulation technique:- In delta modulation, the step size is constant so that its slope overload distortion and granular noise both cannot be controlled. These drawbacks can be controlled by using adaptive delta modulation wherein the step size is variable.	4111
		Thus with adaptive delta modulation the following are the advantages-	
		1. Slope overload distortion and granular noise problem in is reduced.	
		2. Improved signal to noise ratio.	
		3. Wide dynamic range is achieved with variable step size.	
		4.Better bandwidth utilization than delta modulation	
((iii)	Draw and explain spread spectrum modulation system.	4 M
	Ans:	Dram and explain spread speed and mountation system.	2M
1	115.	(Diagram = 2M, Explanation =2M	
		Model Of a Spread Specturm modulation System:-	



	the presence of interference with high probability. • the binary data system 10110010 draw Unipolar RZ, Polar RZ, split phase nchester and AMI.	4M 1M each wave-
b) Att	empt any ONE of the following	6M

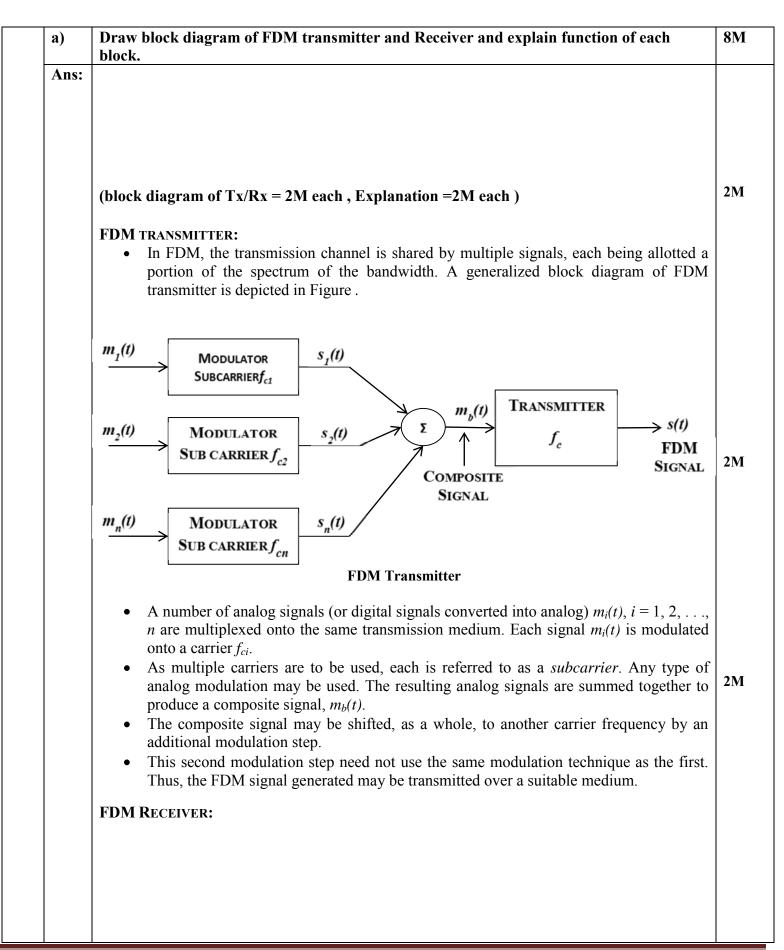


Ans:	· ·	= 2M , Explanation =2M and Example = 2 M) on Of Cyclic Redundancy Check	4
	(CRC):- With CRC sequence of the unit so to binary num At its destin is no remain data unit ha The redund predetermin	the entire data stream is treated as long continuous binary number. In this method, a f redundant bits, called the CRC or the CRC remainder, is appended to the end of that the resulting data unit becomes exactly divisible by a second, predetermined	2
	Example: 4	Data CRC Divisor Data Divisor Data CRC Divisor Remainder CRC Zero, accept Nonzero, reject Nonzero, reject Sender Receiver Sender	
		1 1 1 1 0 1 quotient	
	1 1 0 1 divisor	$ \begin{array}{c} 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 & (data plus extra 3 zero) \\ 1 & 1 & 0 & 1 \\ \hline \hline \hline \hline 1 & 0 & 0 & 0 \\ 1 & 1 & 0 & 1 \\ \hline \hline \hline 1 & 1 & 0 & 1 \\ \hline \hline \hline 0 & 1 & 1 & 0 & 1 \\ \hline \hline 0 & 0 & 1 & 1 & 0 \\ \hline \hline 1 & 1 & 0 & 0 \\ \hline \hline 1 & 1 & 0 & 0 \\ \hline \hline 0 & 0 & 1 & 0 \\ \hline \hline 0 & 0 & 1 & 1 \\ \hline \hline 0 & 0 & 1 & 1 \\ \hline \hline \hline 0 & 0 & 1 & remainder \end{array} $	
	At Rx		

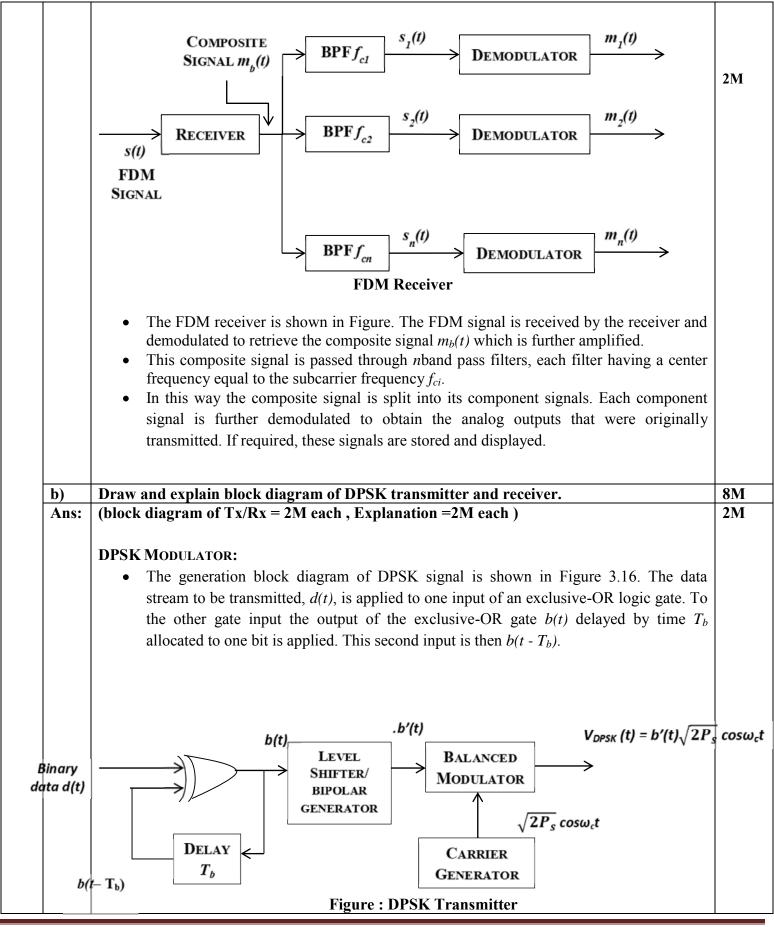


		1 1 1 1 0 1 quotient		
	1 1 0 1 divisor	1 0 0 1 0 0 0 0 1 (data plus CH	RC received)	
		1 0 0 0 1 1 0 1		
		1 0 1 0 1 1 0 1		
		1 1 1 0 1 1 0 1		
		0 1 1 0 0 0 0 0		
		$\begin{array}{c} 1 & 1 & 0 & 1 \\ 1 & 1 & 0 & 1 \end{array}$		
		0 0 0 result		
(ii)	Explain freq	uency hopping. Compare slow a	and fast frequency hopping (four points)	6M
			MFSK. It is the process of modifying frequency	V
	Different type FHSS(Frequence)	nal using frequency hope generate es of frequency hopping:-Slow and ency Hoping Spread spectrum) : adcasting over seemingly random	ed by bits of PN sequence. d fast frequency hopping.	,
	Different type FHSS(Frequencies in frequencies in	es of frequency hopping:-Slow and ency Hoping Spread spectrum): adcasting over seemingly random a	d by bits of PN sequence. d fast frequency hopping.	e
	Different type FHSS(Frequ Signal is broa frequencies in frequency aff	es of frequency hopping:-Slow and ency Hoping Spread spectrum): adcasting over seemingly random so a sync with transmitter. Eavesdrop fects only a few bits.	d by bits of PN sequence. d fast frequency hopping. series of frequencies Receiver hops between opers hear unintelligible blips. Jamming on on	e 1M
	Different type FHSS(Frequencies in frequencies in frequency aff	es of frequency hopping:-Slow and ency Hoping Spread spectrum): idcasting over seemingly random a n sync with transmitter. Eavesdrop	ed by bits of PN sequence. d fast frequency hopping. series of frequencies Receiver hops between	e 1M each
	Different type FHSS(Frequencies in frequencies in frequency aff 1.More to per frequency aff 2.chip ratio	es of frequency hopping:-Slow and ency Hoping Spread spectrum): adcasting over seemingly random s in sync with transmitter. Eavesdrop fects only a few bits. Slow frequency hopping than one symbols are transmitted uency hop ate is equal to symbol rate	d by bits of PN sequence. d fast frequency hopping. series of frequencies Receiver hops between opers hear unintelligible blips. Jamming on on Fast frequency hopping 1.More than one frequency hops are required to transmit one symbol 2.chip rate is equal to hop rate	e 1M each
	Different type FHSS(Frequencies in frequencies in frequency aff 1.More to per frequency 2.chip ra 3.symbol	es of frequency hopping:-Slow and ency Hoping Spread spectrum): adcasting over seemingly random s in sync with transmitter. Eavesdrop fects only a few bits. Slow frequency hopping than one symbols are transmitted uency hop ate is equal to symbol rate of rate is higher than hop rate	d by bits of PN sequence. d fast frequency hopping. series of frequencies Receiver hops between opers hear unintelligible blips. Jamming on on Fast frequency hopping 1.More than one frequency hops are required to transmit one symbol 2.chip rate is equal to hop rate 3.hop rate is higher than symbol rate	e
	Different type FHSS(Frequencies in frequencies in frequency aff 1.More to per frequency 2.chip ra 3.symbol 4.same of	es of frequency hopping:-Slow and ency Hoping Spread spectrum): adcasting over seemingly random son sync with transmitter. Eavesdrop fects only a few bits. Slow frequency hopping than one symbols are transmitted uency hop ate is equal to symbol rate of rate is higher than hop rate carrier frequency is used to	d by bits of PN sequence. d fast frequency hopping. series of frequencies Receiver hops between opers hear unintelligible blips. Jamming on on Fast frequency hopping 1.More than one frequency hops are required to transmit one symbol 2.chip rate is equal to hop rate 3.hop rate is higher than symbol rate 4.one symbol is transmitted over	e 1M each
	Different type FHSS(Frequencies in frequencies in frequency aff 1.More to per frequency 2.chip ra 3.symbol 4.same of transmit 5.A jame	es of frequency hopping:-Slow and ency Hoping Spread spectrum): adcasting over seemingly random s in sync with transmitter. Eavesdrop fects only a few bits. Slow frequency hopping than one symbols are transmitted uency hop ate is equal to symbol rate of rate is higher than hop rate	d by bits of PN sequence. d fast frequency hopping. series of frequencies Receiver hops between opers hear unintelligible blips. Jamming on on Fast frequency hopping 1.More than one frequency hops are required to transmit one symbol 2.chip rate is equal to hop rate 3.hop rate is higher than symbol rate	e 1M each
	Different type FHSS(Frequencies in frequencies in frequency aff 1.More to per frequency 2.chip ra 3.symbol 4.same of transmit 5.A jamic carrier frequency	es of frequency hopping:-Slow and ency Hoping Spread spectrum): adcasting over seemingly random in a sync with transmitter. Eavesdrop fects only a few bits. Slow frequency hopping than one symbols are transmitted uency hop ate is equal to symbol rate of rate is higher than hop rate carrier frequency is used to is one or more symbols mer can detect this signal if the	d by bits of PN sequence. d fast frequency hopping. series of frequencies Receiver hops between opers hear unintelligible blips. Jamming on on Fast frequency hopping 1.More than one frequency hops are required to transmit one symbol 2.chip rate is equal to hop rate 3.hop rate is higher than symbol rate 4.one symbol is transmitted over multiple carrier in different hops 5.A jammer cannot detect this signal because one symbol is transmitted using	e 1M ea











<i>d</i> (<i>t</i>)		$b(t - T_b)$			b(t)		
Logic Level	Voltage	Log	gic Level	Vol	tage	Logic Level	1	Voltage
0	-1		0	-	1	0		-1
0	-1		1		1	1		1
1	1		0	-	1	1		1
1	1		1		1	0		-1
From figure 3.1	6, $b(t)$ is given b	y, b(t) =	$= d(t) \oplus b(t)$	$t-T_b$)	•		•	
Input Data d(<i>t</i>)		1	0	1	1	1	0
Delayed inpu	t $b(t-T_b)$		0	1	1	0	1	0
XOR Output	b(t)	0	1	1	0	1	0	0
Output Phase			0°	0°	180°	0°	180°	1809
DPSK input (Receiver)	at	180°	0°	0°	180°	0°	180°	1809
Recovered da	ta stream		1	0	1	1	1	0

• It is observed that when d(t) = 0, $b(t) = b(t - T_b)$ and when d(t) = 1, $b(t) = \overline{b(t - T_b)}$. As seen in Figure 3.16,b(t) is applied to a level shifter which assigns a positive voltage level when b(t) = 1 and a negative voltage level when b(t) = 0. The level shifter output is then applied to a balanced modulator to which a carrier signal $\sqrt{2P_s cos\omega_c t}$ is also applied. The modulator output, which is the transmitted signal is given by,

$$V_{DPSK}(t) = b'(t) \sqrt{2P_s} \cos \omega_c t$$
$$= (\pm 1) \sqrt{2P_s} \cos \omega_c t$$

DPSK DEMODULATOR:

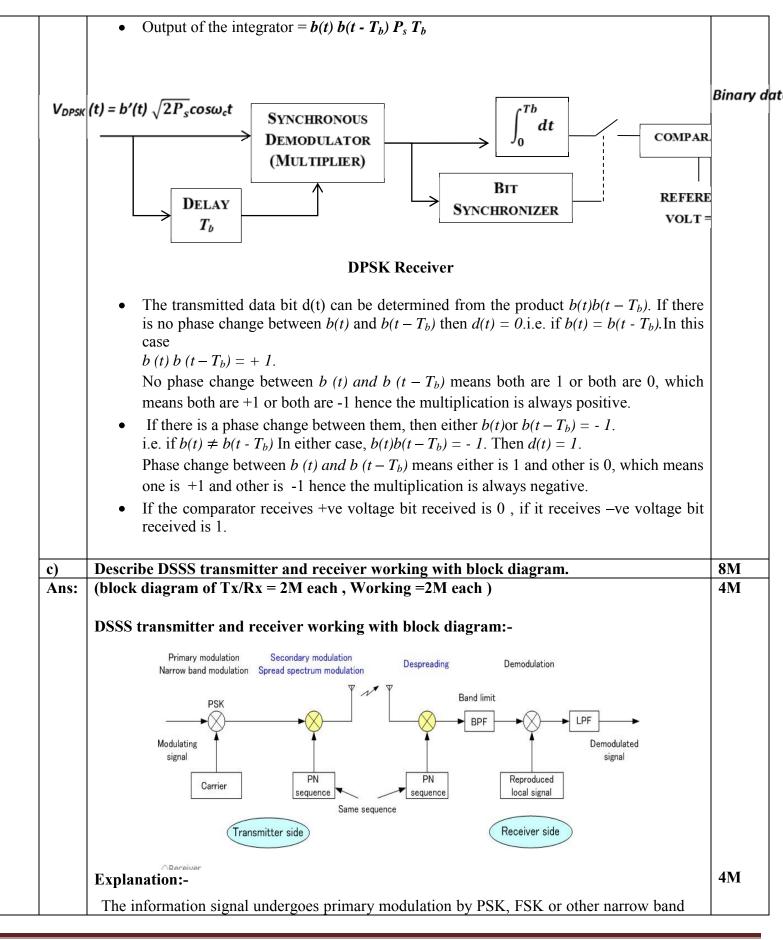
• The DPSK demodulator is shown in Figure. Here the received signal and the received signal delayed by the bit time Tb are applied to a balanced modulator. The balanced modulator output is given by,

 $b(t)b(t-T_b)(2P_s) \cos\omega_c(t) \cos\omega_c(t-T_b)$

• The output of the balanced modulator is applied to the integrator which suppresses the double frequency term. The first term $[b(t)b(t - T_b)P_s \cos\omega_c T_b]$ is the required signal and $\omega_c T_b$ is selected in such a way so that $\omega_c T_b = 2n\pi$ (where *n* is an integer) so that $\cos\omega_c T_b = + 1$ and the signal output will be as large as possible. Further, with this selection, the bit duration encompasses an integral number of clock cycles and the integral of the double frequency term is exactly zero.

2M







modulation and secondary modulation with spread spectrum modulation. Spread spectra are obtained by multiplying the primary modulated signal and the square wave, called the PN sequence. Contrariwise, as with commercial radio, there are cases where spread modulation is applied to the data first, and narrow band modulation such as PSK or FSK is applied afterwards.

The figure below is an example of spread spectrum modulation and demodulation using PSK for primary modulation.

Receiver:

If despreading is applied to the received diffuse wave, it returns to the PSK or FSK modulated wave resulting from primary modulation. Then, as with narrowband demodulation, if the despread wave and local signal are multiplied, and appropriate low pass processing is applied, the information signal is obtained. Despreading involves multiplying the same PN code as that used at the transmitting end for the receiving wave. At this time, it's necessary to synchronize the receiving wave and PN code.

There are two processing methods on the receiving side, demodulation of the information signal after despreading, and obtaining a positive and negative PN code by multiplying the local signal by the receiving wave and despreading using correlation detection. With the former there is process gain but the problem of synchronization remains. With the latter, the spectrum density of the receiving wave itself is low, and regeneration of the local carrier for performing synchronous detection is a problem. Commercial SS radio equipment uses the latter, but it requires considerable power and has a short communication range.

Despreading:

The signal that enters the antenna of the receiver includes outside interference waves and noise. If this signal is despread, the signal component returns to a narrowband modulated wave and the interference components are diffused, expanding the spectrum infinitely so that its power density falls. Therefore, by inputting the signal with frequency band restricted using a BPF, the interference component power that falls into the demodulation frequency band is reduced. The occurrence of errors is calculated using a stochastic process, so ultimately, using a spread spectrum results in fewer errors, and this is why spread spectrum communication is resistant to interference.

Demodulation:

Demodulation is normal narrowband demodulation. The local signal is regenerated from the receiving wave and after multiplication by the receiving wave, unnecessary components are eliminated with an LPF. Primary modulation uses PSK, so synchronous detection is necessary.

PNsequence :

The PN sequence is switched at a far faster speed than the symbol rate of the information signal and its spectrum covers a wide band. For this reason, the spectrum of the modulated wave after primary modulation also covers a wide band. We won't go into detail here, but PN sequences must meet the conditions required for spread spectrum modulation such as the relationship of the numbers 1 and 0.



